



Management Science

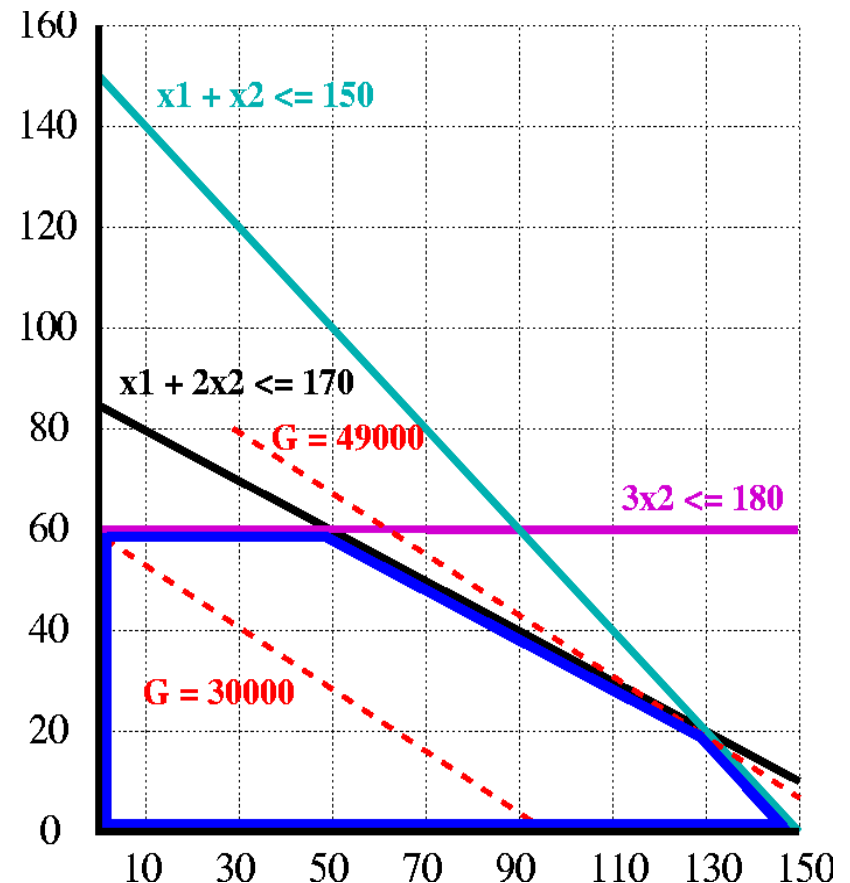
Integer and Mixed

Programming

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LP Problems Today

- Today the main difficulty is formulating the LP problem, not in solving it !!!



Packing Problem



- In this type of problem we have to pack a number of items into a number of physical or logical spaces
- We want to optimize the number of items or the value of the items packed into the spaces
- There are constraints on how much can be packed into the spaces based on characteristics of the “items” and characteristics of the “space”



Capacity of Plane



Space	Weight Capacity (tons)	Volume Capacity (cu ft)
Forward Hold	75	4000
Main Hold	150	10000
Aft Hold	50	7000

Available Cargo

[which should we take none of, which should we take all of]

Cargo	Tons	Cu ft/ton	\$/ton
Memory Chips	200	60	800
NIC's	100	48.6	600
Monitors	500	4.1 (dense)	200 (cheap)
Mother Boards (w cpu)	50	240 (bulky)	2500 (expensive)

- What are the variables ?
- What is the objective function ?
- What are the constraints ?



■ Do not look ahead !



Variables

- Let:
 - W = tons of memory chips
 - X = tons of NIC's
 - Y = tons of monitors
 - Z = tons of mother boards



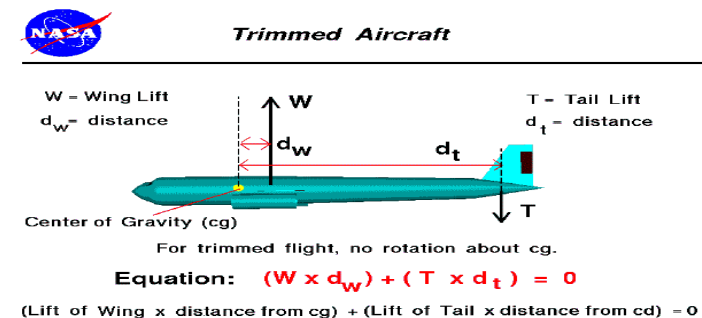
Objective Function

- Objective function to **maximize**:
 - $Z = 800 W + 600 X + 200 Y + 2500 Z$



Constraints

- Availability of cargo (supply)
- Weight of each hold
- Capacity (volume) of each hold
- What if we do not load the holds “evenly” by weight ?

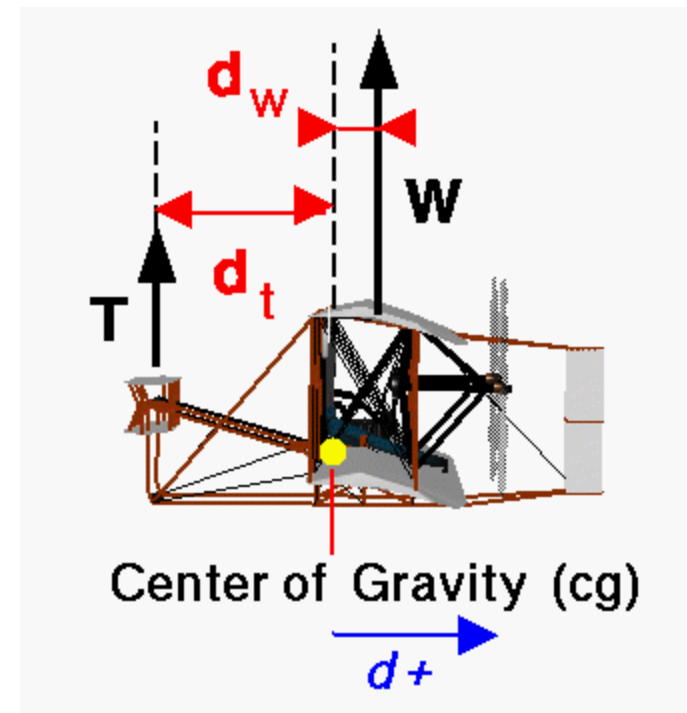


■ Do not look ahead !



Distribution Variables

- More variables:
 - $WF = W$ in forward hold (chips in forward hold)
 - $WC = W$ in main hold
 - $WA = W$ in aft hold
 - Similarly for other items:
 - XF, XC, XA
 - YF, YC, YA
 - ZF, ZC, ZA



Distribution Constraints

- $W = WF + WC + WA$
- $X = XF + XC + XA$
- $Y = YF + YC + YA$
- $Z = ZF + ZC + ZA$

Trim Constraints

- The ratio of the weight of the cargo stored in each hold to that hold's maximum weight must be equal across all three holds !
- $(WF + XF + YF + ZF)/75 = (WC + XC + YC + ZC)/150$
- $(WA + XA + YA + ZA)/50 = (WC + XC + YC + ZC)/150$

Supply Constraints

- $W \leq 200$ (chips)
- $X \leq 100$ (NIC's)
- $Y \leq 500$ (monitors)
- $Z \leq 50$ (boards)

Weight Constraints

- $WF + XF + YF + ZF \leq 75$ (forward)
- $WC + XC + YC + ZC \leq 150$ (center)
- $WA + XA + YA + ZA \leq 50$ (aft)

Volume Constraints

- $60 WF + 48.6 XF + 4.1 YF + 240 ZF \leq 4000$ (forward)
- $60 WC + 48.6 XC + 4.1 YC + 240 ZC \leq 10000$ (center)
- $60 WA + 48.6 XA + 4.1 YA + 240 ZA \leq 7000$ (aft)

Excel Solution

	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U				
1	Packing Problem																								
2																									
3		Total Amount				Loading																			
4	Decision variables:	W	X	Y	Z	WF	WC	WA	XF	XC	XA	YF	YC	YA	ZF	ZC	ZA								
5	units	200	0	38	37	56	144	0	0	0	0	17	0	21	2.5	5.6	29								
6																									
7	Objective function:																	Total Profit							
8	profit per unit	800	600	200	2500													\$259,750.95							
9																									
10	Constraints:																	Lhs	Type	Rhs	Units				
11	Define W	-1				1	1	1										-2.84217E-14	=	0	tons				
12	Define X		1						-1	-1	-1							0	=	0	tons				
13	Define Y			1								-1	-1	-1				-7.10543E-15	=	0	tons				
14	Define Z				1										-1	-1	-1	-7.10543E-15	=	0	tons				
15	Trim 1					2	-1		2	-1		2	-1		2	-1		1.15463E-14	=	0	---				
16	Trim 2						1	-3		1	-3		1	-3		1	-3	-1.42109E-14	=	0	---				
17	Forward Hold Limit					1			1			1			1			75	<=	75	tons				
18	Main Hold Limit						1			1			1			1		150	<=	150	tons				
19	Aft Hold Limit							1			1			1			1	50	<=	50	tons				
20	W supply	1																200	<=	200	tons				
21	X supply		1															0	<=	100	tons				
22	Y supply			1														38.15175922	<=	500	tons				
23	Z supply				1													36.84824078	<=	50	tons				
24	Fwd vol limit					60			48.6			4.1			240			4000	<=	4000	tons				
25	Cnt vol limit						60			48.6			4.1			240		10000	<=	10000	tons				
26	Aft vol limit							60			48.6			4.1			240	7000	<=	7000	tons				
27																									

Solver

[no NIC's, all the memory chips]

	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U
1	Packing Problem																				
2																					
3		Total Amount						Loading													
4	Decision variables:	W	X	Y	Z	WF	WC	WA	XF	XC	XA	YF	YC	YA	ZF	ZC	ZA				
5	units	200	0	38	37	56	144	0	0	0	0	17	0	21	2.5	5.6	29				
6																					
7	Objective function:																				
8	profit per unit	800	600	200	2500													Total Profit			
9																		\$259,750.95			
10	Constraints:																	Lhs	Type	Rhs	Units
11	Define W	-1				1	1	1										-2.84217E-14	=	0	tons
12	Define X		1						-1	-1	-1							0	=	0	tons
13	Define Y			1																	
14	Define Z				1																
15	Trim 1					2	-1														
16	Trim 2						1														
17	Forward Hold Limit					1															
18	Main Hold Limit						1														
19	Aft Hold Limit																				
20	W supply	1																			
21	X supply		1																		
22	Y supply			1																	
23	Z supply				1																
24	Fwd vol limit					60															
25	Cnt vol limit						60														
26	Aft vol limit																				
27																					
28																					
29																					
30																					
31																					

Set Target Cell:	\$R\$8	Solve
Equal To:	<input checked="" type="radio"/> Max <input type="radio"/> Min <input type="radio"/> Value of:	0
By Changing Cells:	\$B\$5:\$Q\$5	Guess
Subject to the Constraints:	\$R\$11:\$R\$16 = \$T\$11:\$T\$16 \$R\$17:\$R\$26 <= \$T\$17:\$T\$26	Add Change Delete
		Options
		Reset All
		Help

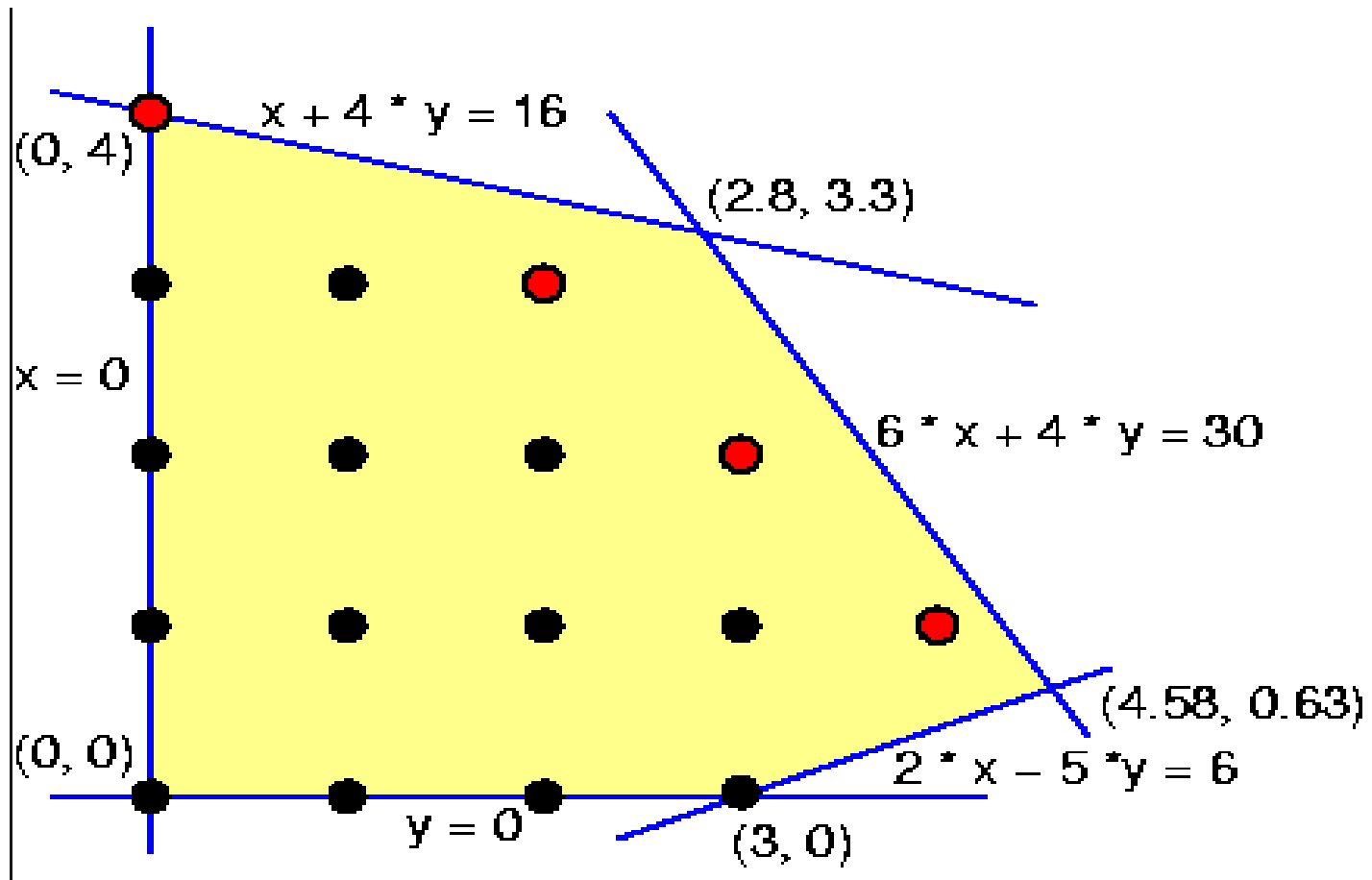
Integer Programming

- An integer programming model is one where one or more of the decision variables has to take on an integer value in the final solution
- There are three types of integer programming problems
 1. **Pure integer programming** where all variables have integer values
 2. **Mixed-integer programming** where some but not all of the variables will have integer values
 3. Zero-one integer (**binary programming**) are special cases in which all the decision variables must have integer solution values of 0 or 1

Integer Programming (con't)

- Solving an integer programming problem is much more difficult than solving an LP problem
- Even the fastest computers can take an excessively long time to solve big integer programming problems
- There are a number of methods for solving integer programming problems, and a common technique used to solve integer programming problems is the branch and bound method
- There may be multiple solutions to integer programming problems

Integer Programming (con't)



Example of Integer Programming

- A company produces two products popular with home renovators: old-fashioned chandeliers and ceiling fans
- Both the chandeliers and fans require a two-step production process involving wiring and assembly
- It takes about 2 hours to wire each chandelier and 3 hours to wire a ceiling fan
- Final assembly of the chandeliers and fans requires 6 and 5 hours respectively
- The production capability is such that only 12 hours of wiring time and 30 hours of assembly time are available

Example of Integer Programming (con't)

- Each chandelier produced nets the firm \$7 and each fan \$6
- What are the variables ?
- What are the constraints ?



■ Do not look ahead !



Example of Integer Programming (con't)

- This “production mix” decision can be formulated using LP as follows:

$$\begin{array}{ll} \text{Maximize profit} = & \$7X_1 + \$6X_2 \\ \text{subject to} & 2X_1 + 3X_2 \leq 12 \quad (\text{wiring hours}) \\ & 6X_1 + 5X_2 \leq 30 \quad (\text{assembly hours}) \\ & X_1, X_2 \geq 0 \quad (\text{nonnegative}) \end{array}$$

where

X_1 = number of chandeliers produced, **must be an integer**
 X_2 = number of ceiling fans produced, **must be an integer**



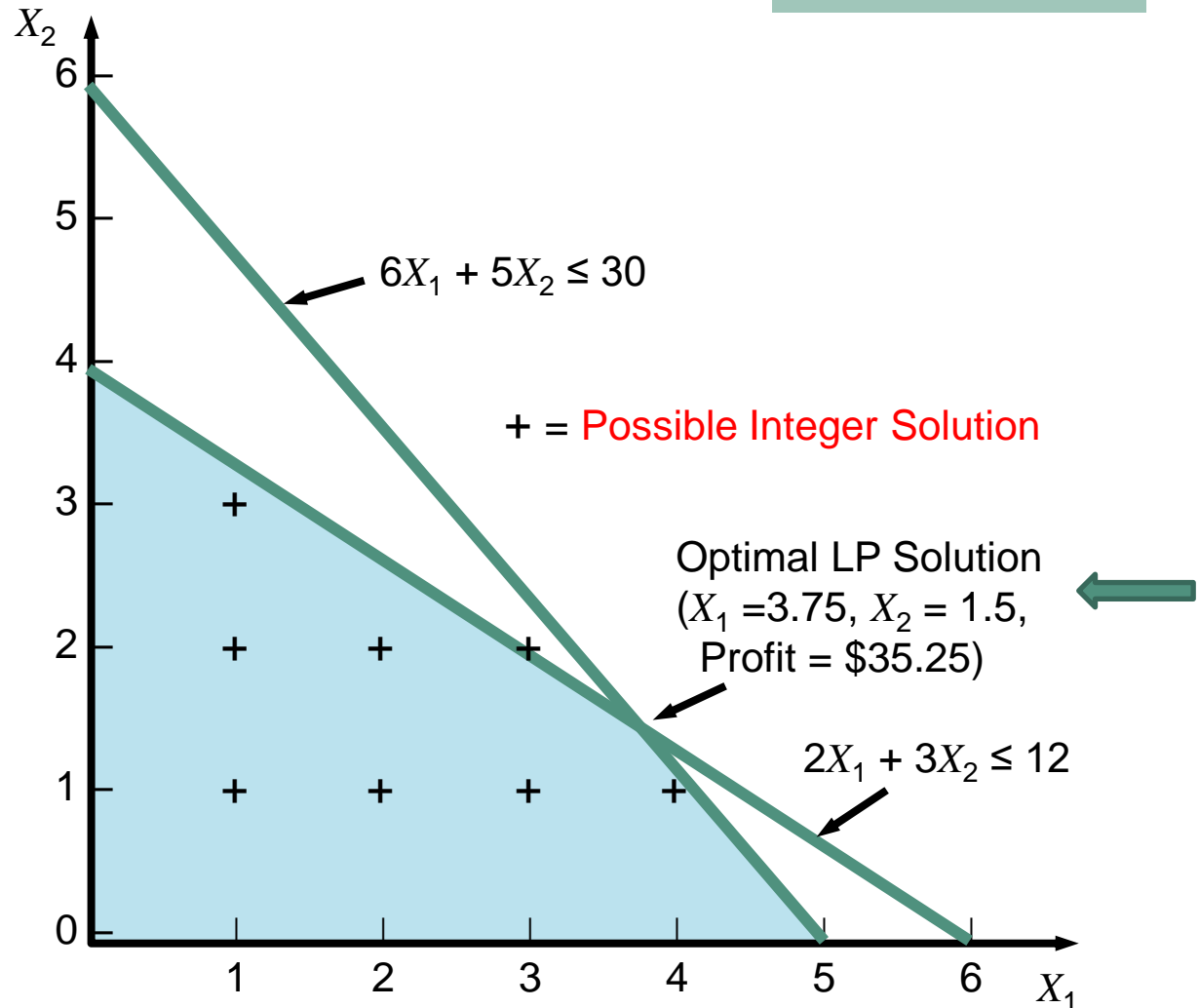
- How could we approach solving this problem ?



■ Do not look ahead !



Example of Integer Programming (con't)



Example of Integer Programming (con't)

- A first attempt at solving may be to round the optimal values to $X_1 = 4$ and $X_2 = 2$
- However, this is outside of the feasible area
- Rounding X_2 down to 1 gives a feasible solution, but it may not be *optimal*
- This could be solved using an *enumeration* method
- Enumeration is generally not feasible for large problems

Enumeration For Integer Programming

Integer solutions

CHANDELIERS (X_1)	CEILING FANS (X_2)	PROFIT ($\$7X_1 + \$6X_2$)
0	0	\$0
1	0	7
2	0	14
3	0	21
4	0	28
5	0	35
0	1	6
1	1	13
2	1	20
3	1	27
4	1	34
0	2	12
1	2	19
2	2	26
3	2	33
0	3	18
1	3	25
0	4	24

← Optimal solution to integer programming problem

Example of Integer Programming (con't)

- The optimal integer solution of $X_1 = 5$, $X_2 = 0$ gives a profit of \$35
- The optimal integer solution (35) is less than the optimal LP solution (35.25)
- An integer solution can *never* be better than the LP solution and is *usually* a lesser solution

Example of Integer Programming (con't)

	A	B	C	D	E	F	G	
1								
2			X1	X2				
3		Solution	0	0				
4		Con 1	2	3	=C3*C4 + D3*D4	12	<=	
5		Con 2	6	5	=C3*C5 + D3*D5	30	<=	
6		Objective	7	6	=C3*C6 + D3*D6		Min MAX	
7								
8								
9								

Example of Integer Programming (con't)

	A	B	C	D	E	F	G	H	I	J	K	L	M
1													
2			X1	X2									
3		Solution	0	0									
4		Con 1	2	3	0	12	<=						
5		Con 2	6	5	0	30	<=						
6		Objective	7	6	0								
7													
8													
9													
10													
11													
12													
13													
14													
15													
16													
17													
18													
19													
20													
21													

Max

Solver Parameters

Set Target Cell:

Equal To: ☒ Max ☐ Min ☐ Value of:

By Changing Cells:

Subject to the Constraints:

For integer problems, one still chooses the “simplex” method in Excel, Excel will automatically switch over to the “branch and bound” method

Example of Integer Programming (con't)

	A	B	C	D	E	F	G	H	I	J	K	L	M
1													
2			X1	X2									
3		Solution	5	0									
4		Con 1	2	3	10	12	<=						
5		Con 2	6	5	30	30	<=						
6		Objective	7	6	35		Min						
7													
8													
9													
10													
11													
12													
13													
14													
15													
16													
17													
18													
19													

Solver Results

Solver found a solution. All constraints and optimality conditions are satisfied.

☒ Keep Solver Solution
 ☐ Restore Original Values

Reports

Answer
Sensitivity
Limits

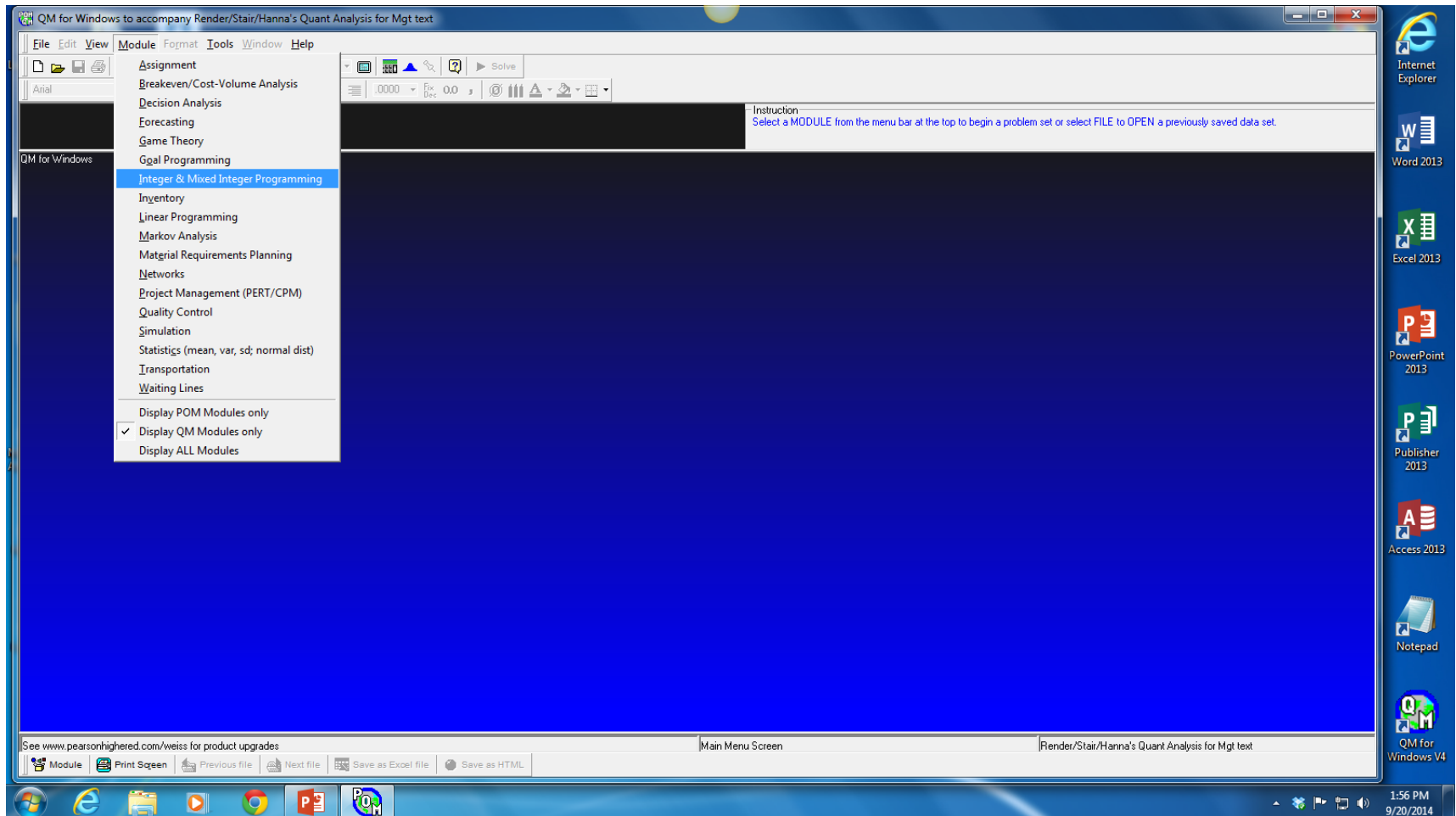
OK

Cancel

Save Scenario...

Help

Using QM...



Using QM...

Create data set for Integer && Mixed Integer Programming

Title:

Number of Constraints

Number of Variables

Objective
☒ Maximize
☐ Minimize

Row names Column names Overview

☒ Constraint 1, Constraint 2, Constraint 3, ...
☐ a, b, c, d, e, ...
☐ A, B, C, D, E, ...
☐ 1, 2, 3, 4, 5, ...
☐ January, February, March, April, ...

☐ Other

QM for Windows - [Data Table]

File Edit View Module Format Tools Window Help

File Edit View Module Format Tools Window Help

Objective

Maximize

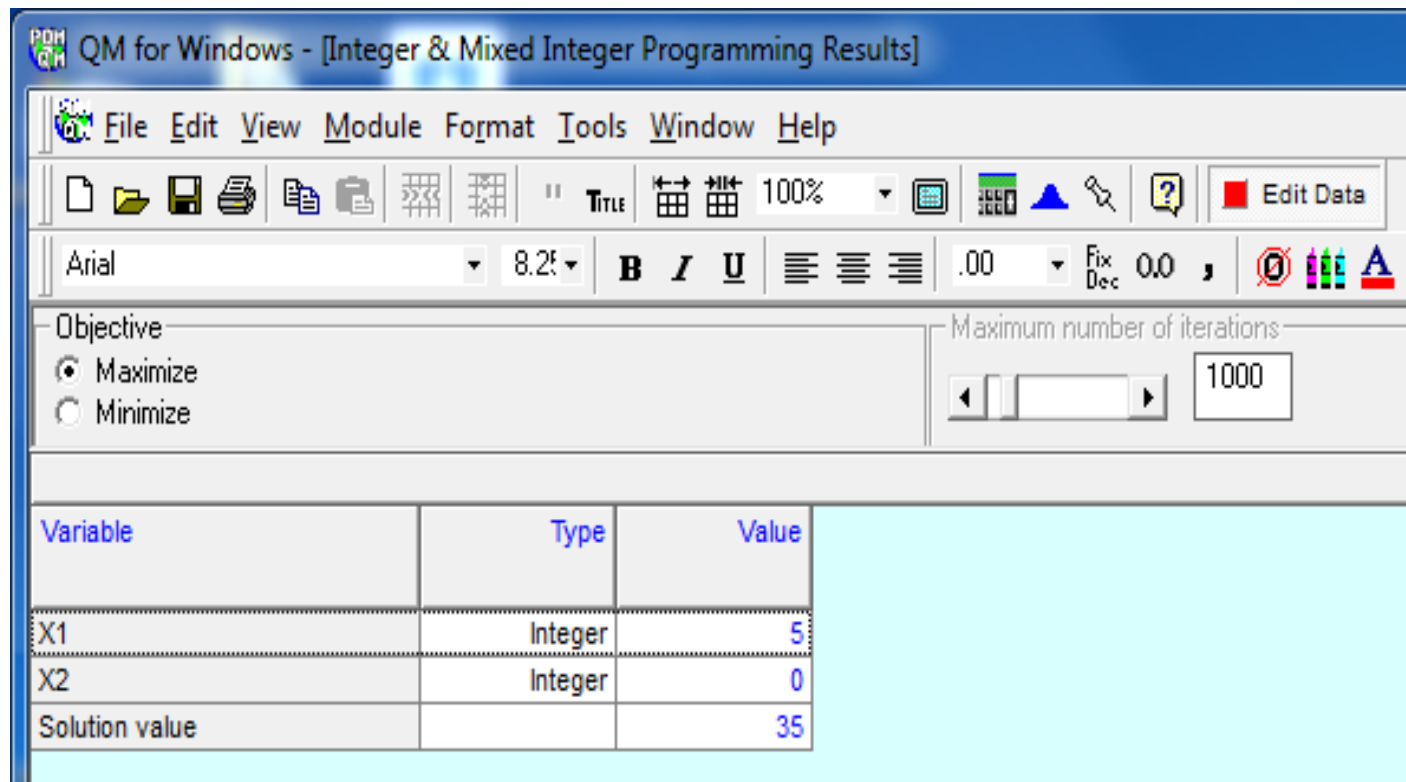
Minimize

Maximum number of iterations

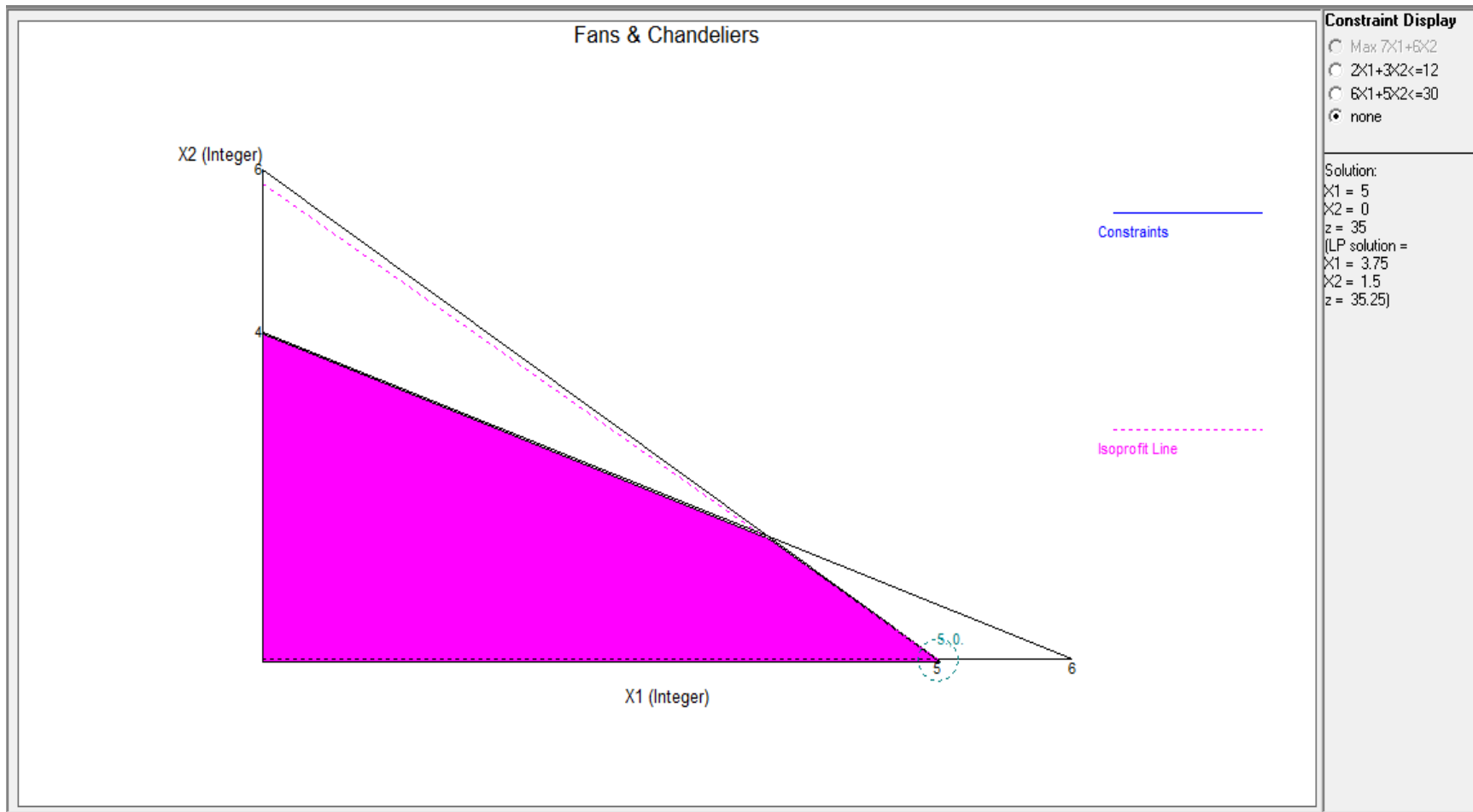
1000

	X1	X2		RHS	Equation form
Maximize	7	6			Max $7X_1 + 6X_2$
Constraint 1	2	3	\leq	12	$2X_1 + 3X_2 \leq 12$
Constraint 2	6	5	\leq	30	$6X_1 + 5X_2 \leq 30$
Variable type	Integer	Integer			

Using QM...



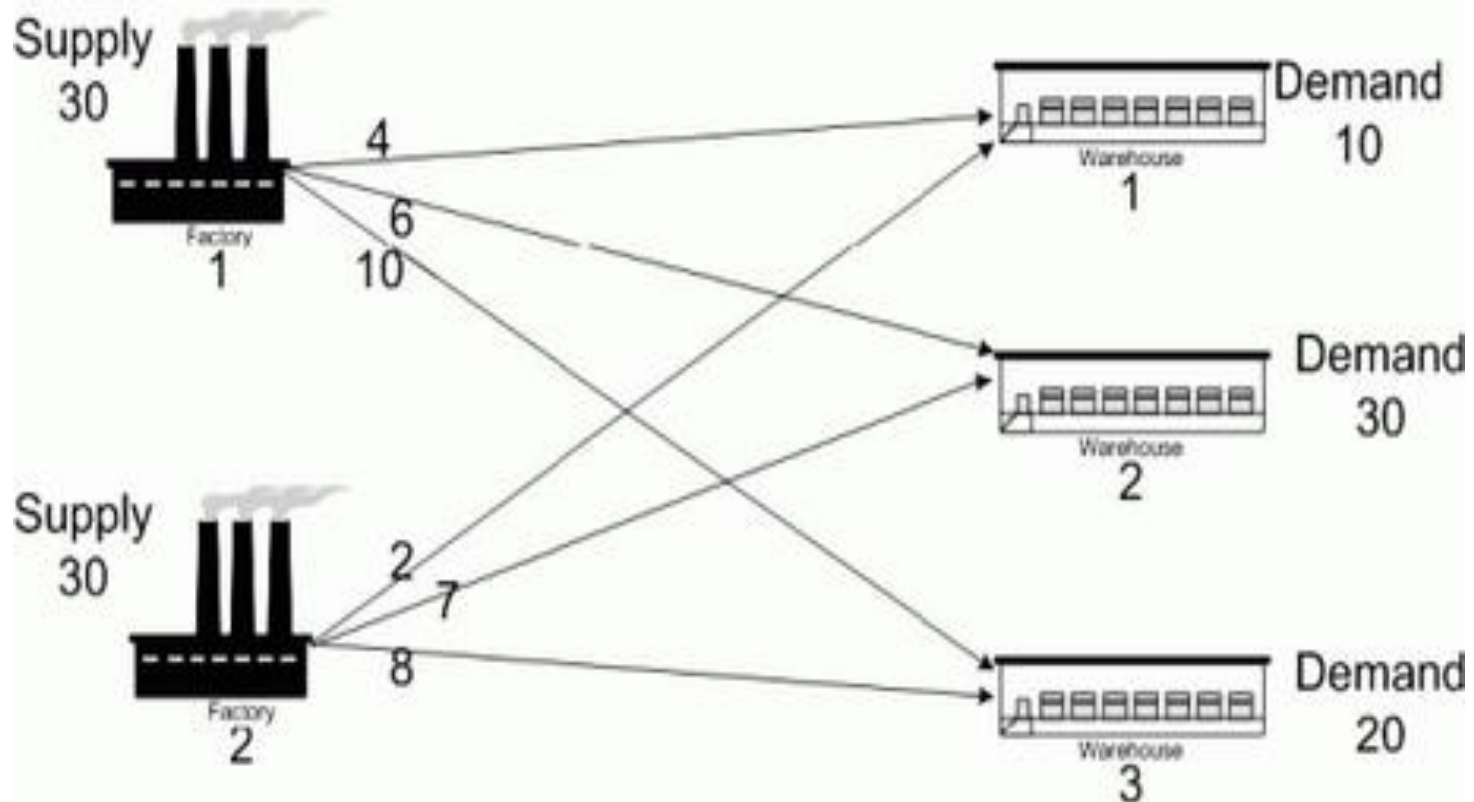
Using QM...



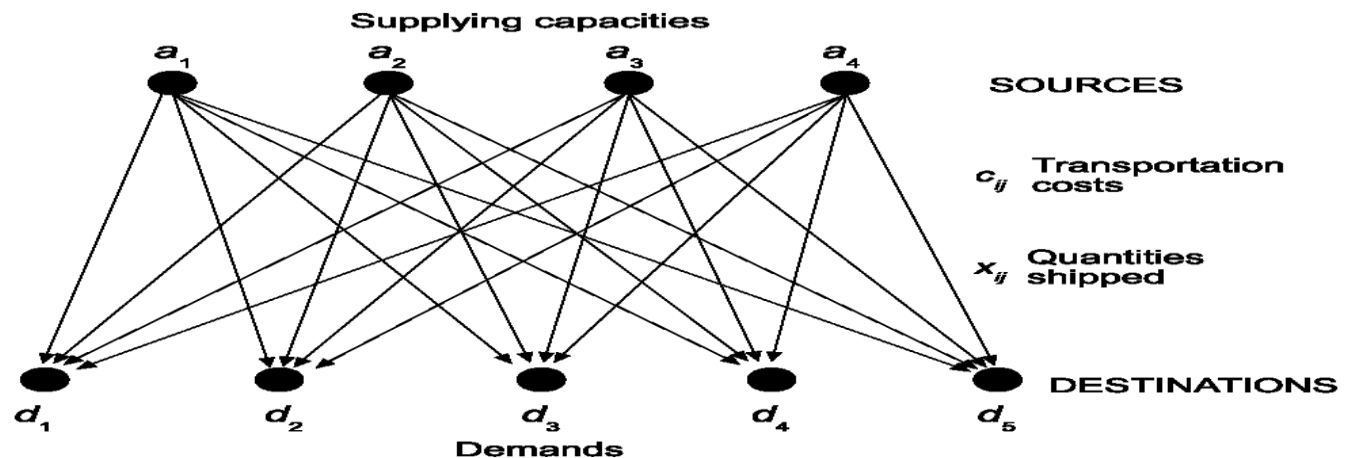
“Transportation” Problems

- There are a number of special management problems that can be solved via integer linear programming
- One common problem is that of the “transportation problem”
- Here we have a number of plants (or warehouses) that can produce (or ship) products subject to **supply constraints** and we have a number of consumers (or customers) that have certain **demands for products**
- There is a **cost associated with shipping product from the warehouses to the customers**

“Transportation” Problems (con’t)



- What are the variables ?
- What is the objective function ?
- What are the constraints ?



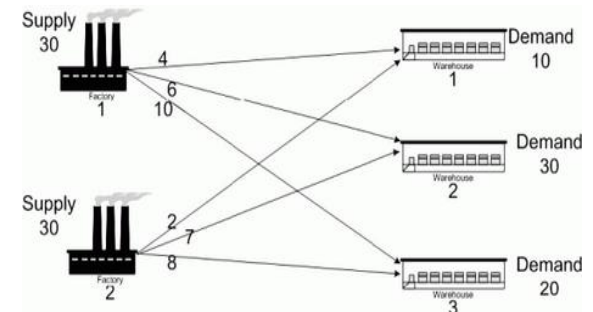
■ Do not look ahead !



“Transportation” Problems (con’t)

■ Let:

- X_{ij} = number of units shipped from source i to destination j (*integers*)
- c_{ij} = cost of one unit from source i to destination j
- s_i = supply at source i
- d_j = demand at destination j



“Transportation” Problems (con’t)

Minimize cost =

$$\sum_{j=1}^n \sum_{i=1}^m c_{ij} x_{ij}$$

Subject to:

$$\sum_{j=1}^n x_{ij} \leq s_i \quad i = 1, 2, \dots, m \quad [\text{supply}]$$

$$\sum_{i=1}^m x_{ij} \leq d_j \quad j = 1, 2, \dots, n \quad [\text{demand}]$$

$x_{ij} \geq 0$ for all i and j , *and are integers*

Transportation Matrix in Excel

[showing demand, supply, and transport costs]

Microsoft Excel - transportation.xls

File Edit View Insert Format Tools Data Window Help

C14 fx

	A	B	C	D	E	F	G
1							
2							
3			Customers				
4		FROM\TO	Birmingham	New Orleans	Jackson	SUPPLY	
5	Warehouses	Memphis	5	4	3	100	
6		Dallas	8	4	3	300	
7		Atlanta	9	7	5	300	
8		DEMAND	300	200	200	700/700	
9							
10							

Cost from
Dallas to
Jackson

What are the cells for the decision variables ?

Set Up Shipment Table and Cost Function

Microsoft Excel - transportation.xls

File Edit View Insert Format Tools Data Window Help

D20 $\text{=SUMPRODUCT}(C6:E8,C15:E17)$

	A	B	C	D	E	F	G
2	Problem - Transportation Parameters						
3							
4			Customers				
5		FROM\TO	Birmingham	New Orleans	Jackson	SUPPLY	
6	Warehouses	Memphis	5	4	3	100	
7		Dallas	8	4	3	300	
8		Atlanta	9	7	5	300	
9		DEMAND	300	200	200	700/700	
10							
11	Solution - Shipments from Warehouses to Customers						
12							
13			Customers				
14		FROM\TO	Birmingham	New Orleans	Jackson	ROW TOTAL	
15	Warehouses	Memphis	0	0	0	0	
16		Dallas	0	0	0	0	
17		Atlanta	0	0	0	0	
18		COL TOTAL	0	0	0	0/0	
19							
20		TOTAL COST ->					
21							

- Target is D20 (minimize)
- Manipulate c15:e17
- Subject to:
 - Col totals equal demand
 - Row totals \leq supply
 - Shipments are integers and +

Need another matrix for decision variables.

Parameters for Excel Solver

Microsoft Excel - transportation.xls

File Edit View Insert Format Tools Data Window Help

Type a question for help

D20 fx =SUMPRODUCT(C6:E8,C15:E17)

	A	B	C	D	E	F	G	H	I
2	Problem - Transportation P								
4				Customers					
5		FROM\TO	Birmingham	New Orleans					
6	Warehouses	Memphis	5	4					
7		Dallas	8	4					
8		Atlanta	9	7					
9		DEMAND	300	200					
10									
11	Solution - Shipments from Warehouse								
12				Customers					
14		FROM\TO	Birmingham	New Orleans					
15	Warehouses	Memphis	0	0	0	0			
16		Dallas	0	0	0	0			
17		Atlanta	0	0	0	0			
18		COL TOTAL	0	0	0	0/0			
19									
20	TOTAL COST ->			0					
21									

Solver Parameters

Set Target Cell: $\text{\$D\$20}$

Equal To: ☐ Max ☒ Min ☐ Value of: 0

By Changing Cells: $\text{\$C\$15:\$E\$17}$

Subject to the Constraints:

- $\text{\$C\$15:\$E\$17} = \text{integer}$
- $\text{\$C\$15:\$E\$17} \geq 0$
- $\text{\$C\$18:\$E\$18} = \text{\$C\$9:\$E\$9}$
- $\text{\$F\$15:\$F\$17} \leq \text{\$F\$6:\$F\$8}$

Buttons: Solve, Close, Options, Add, Change, Delete, Reset All, Help

Excel Solution

Microsoft Excel - transportation.xls

File Edit View Insert Format Tools Data Window Help

Type a question for help

Formula Bar: $\text{=SUMPRODUCT}(C6:E8,C15:E17)$

	A	B	C	D	E	F	G	H	I
2									
3									
4									
5		FROM\TO	Birmingham	New Orleans	Jackson				
6	Warehouses	Memphis	5	4					
7		Dallas	8	4					
8		Atlanta	9	7					
9		DEMAND	300	200					
10									
11									
12									
13									
14		FROM\TO	Birmingham	New Orleans	Jackson	ROW TOTAL			
15	Warehouses	Memphis	99.99999933	3.33333E-07	3.33333E-07	100			
16		Dallas	0	199.9999997	100.0000003	300			
17		Atlanta	200.0000007	0	99.99999933	300			
18		COL TOTAL	300	200	200	700/700			
19									
20		TOTAL COST ->		3900.000001					
21									

Solver Results

Solver found a solution. All constraints and optimality conditions are satisfied.

Reports: Answer, Sensitivity, Limits

☒ Keep Solver Solution
☐ Restore Original Values

OK Cancel Save Scenario... Help

After Formatting Cells

Microsoft Excel - transportation.xls

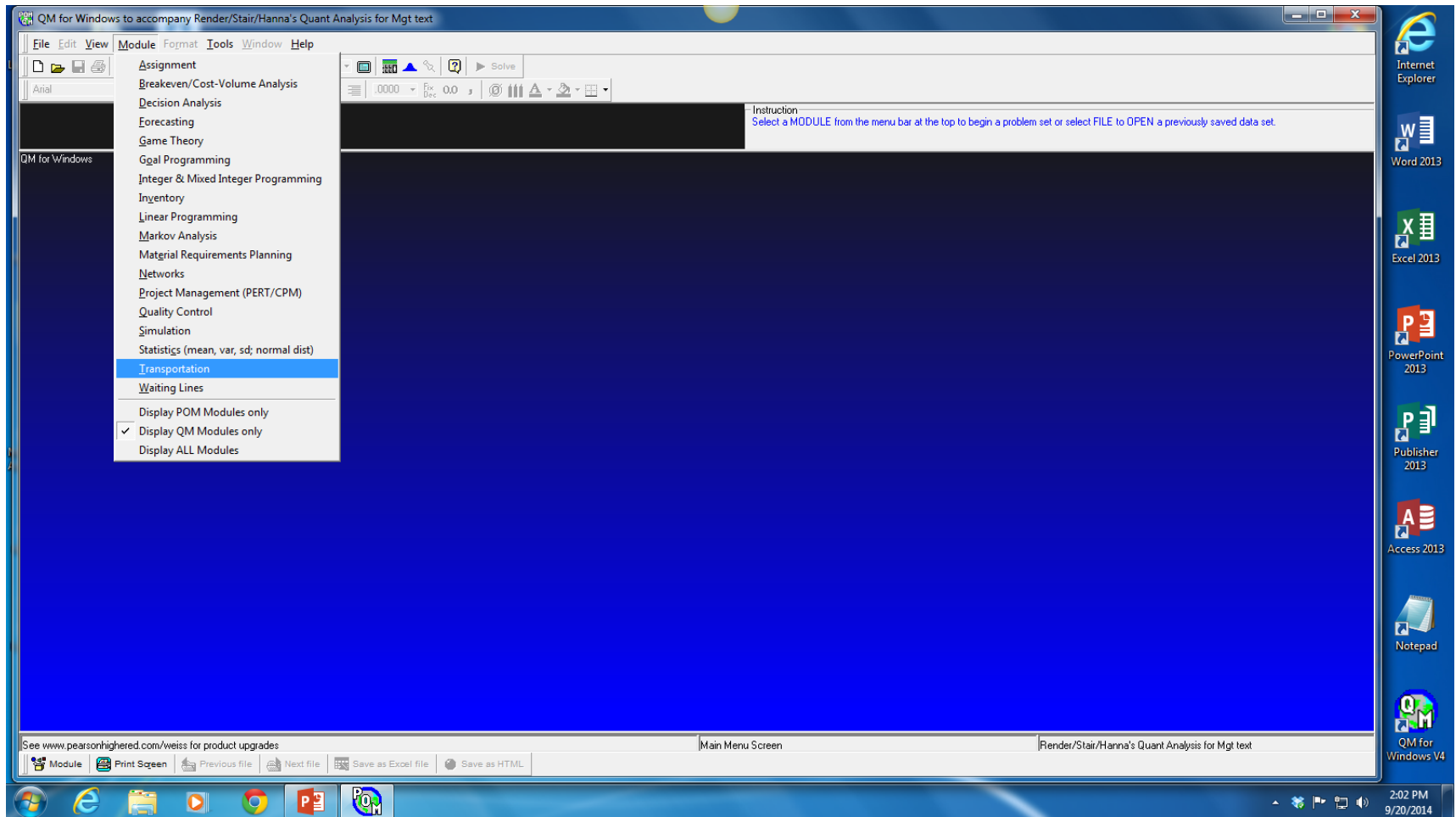
File Edit View Insert Format Tools Data Window Help

Type a question for help

D20 fx =SUMPRODUCT(C6:E8,C15:E17)

	A	B	C	D	E	F	G	H	I
2	Problem - Transportation Parameters								
4			Customers						
5		FROM\TO	Birmingham	New Orleans	Jackson	SUPPLY			
6	Warehouses	Memphis	5	4	3	100			
7		Dallas	8	4	3	300			
8		Atlanta	9	7	5	300			
9		DEMAND	300	200	200	700/700			
10									
11	Solution - Shipments from Warehouses to Customers								
12									
13			Customers						
14		FROM\TO	Birmingham	New Orleans	Jackson	ROW TOTAL			
15	Warehouses	Memphis	100	0	0	100			
16		Dallas	0	200	100	300			
17		Atlanta	200	0	100	300			
18		COL TOTAL	300	200	200	700/700			
19									
20		TOTAL COST ->	\$3,900.00						

Using QM...



Using QM...

Create data set for Transportation

Title: Modify default title

Number of Origins:

Number of Destinations:

Objective:

- ☐ Maximize
- ☒ Minimize

Row names: Column names Overview

- ☒ Origin 1, Origin 2, Origin 3,...
- ☐ a, b, c, d, e, ...
- ☐ A, B, C, D, E, ...
- ☐ 1, 2, 3, 4, 5, ...
- ☐ January, February, March, April, ...
- ☐ Other

Cancel Help OK

Using QM...

The screenshot displays the 'QM for Windows - [Data Table]' application window. The interface includes a menu bar (File, Edit, View, Module, Format, Tools, Window, Help) and a toolbar with various icons for file operations, formatting, and solving. Below the toolbar, there are settings for font (Arial, 8.25), bold, italic, underline, and alignment. The 'Objective' section has radio buttons for 'Maximize' and 'Minimize', with 'Minimize' selected. The 'Starting method' section has a text box containing 'Any starting method'. The main area contains a data table for a transportation problem.

	Birmingham	New Orleans	Jackson	SUPPLY
Memphis	5	4	3	100
Dallas	8	4	3	300
Atlanta	9	7	5	300
DEMAND	300	200	200	

Using QM...

The screenshot displays the 'QM for Windows' application window. The menu bar includes File, Edit, View, Module, Format, Tools, Window, and Help. The toolbar contains various icons for file operations, formatting, and solving. The 'Objective' section has radio buttons for 'Maximize' and 'Minimize', with 'Minimize' selected. The 'Starting method' dropdown is set to 'Any starting method'. The main area is titled 'Transportation Shipments' and contains a table with the following data:

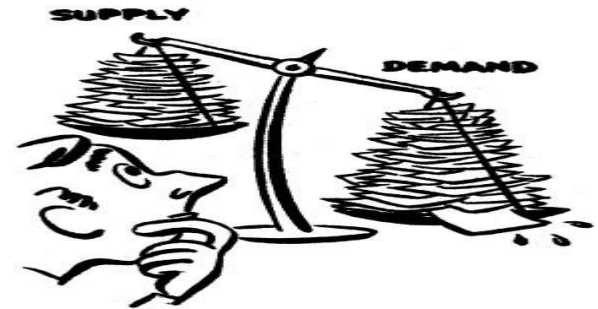
	Birmingham	New Orleans	Jackson
Optimal cost = \$3900			
Memphis	100		
Dallas		200	100
Atlanta	200		100

Unbalanced Transportation Problems

- In real-life problems, **total demand may not be equal to total supply**
- These *unbalanced problems* can be handled easily by introducing *dummy sources* or *dummy destinations*
- If total supply is greater than total demand, a dummy destination (warehouse), with demand exactly equal to the surplus, is created
- If total demand is greater than total supply, we introduce a dummy source (factory) with a supply equal to the excess of demand over supply

Unbalanced Transportation Problems (con't)

- In either case, shipping cost coefficients of zero are assigned to each dummy location or route as no goods will actually be shipped.
- Any units assigned to a dummy destination represent excess capacity
- Any units assigned to a dummy source represent unmet demand



More Than One Optimal Solution

- It is possible for a transportation problem to have **multiple optimal solutions**
- This means that it is possible to design alternative shipping routes with the same total shipping cost
- In the real world, alternate optimal solutions provide management with greater flexibility in selecting and using resources

Unacceptable Or Prohibited Routes

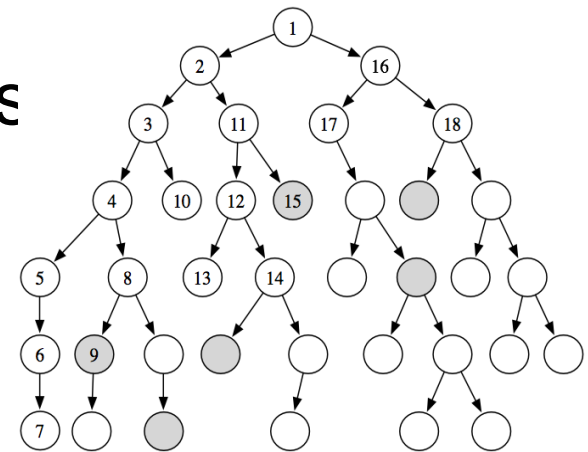
- At times there are transportation problems in which **one of the sources is unable to ship to one or more of the destinations**
 - The problem is said to have an *unacceptable* or *prohibited route*
- In a minimization problem, such a prohibited route is assigned a very high cost to prevent this route from ever being used in the optimal solution
- In a maximization problem, the very high cost used in minimization problems is given a negative sign, turning it into a very bad profit

Facility Location Analysis

- The transportation method is especially useful in helping a firm to decide **where to locate** a new factory or warehouse
- Each alternative location should be analyzed within the framework of one **overall** distribution system
- The new location that yields the minimum cost for the **entire system** is the one that should be chosen

Special Methods

- Transportation and assignment problems have some special algorithms that can be utilized as well as the general purpose integer linear programming methods including:
 - **Branch and bound**
 - Modified corner point methods
 - Stepping stone methods
 - Hungarian methods, etc.



Labor Planning

- Labor planning problems address resource needs over time and/or space
- The resources are often restricted to integers such as X number of people
- For example a bank may need a certain number of tellers for time periods of the day to satisfy differing customer demand by time period as shown on the next slide (for a total of 112 needed daily hours)

Labor Planning (con't)

Time Period	# of Tellers Required
9am – 10am	10
10am – 11am	12
11am – Noon	14
Noon – 1pm	16
1pm – 2pm	18
2pm – 3pm	17
3pm – 4pm	15
4pm – 5pm	10

Labor Planning (con't)

- The bank employs up to 12 full time tellers at a cost of \$100 per day per teller
- Full time tellers work from 9am until 5pm with 1 hour off for lunch (35 hour work week); half of the full time tellers go to lunch at 11am and the other half at noon
- Part time tellers put in exactly 4 hours per day and can start at any hourly even time slot (between 9am and 1pm); they cost \$32 per day (Let P1 be tellers at first time slot, P2 at second time slot, etc.)
- By regulations part time worker hours cannot exceed 50% of the total teller required hours:
 - $4 * (P1 + P2 + P3 + P4 + P5) \leq .5 * 112 \leq 56$

Labor Planning (con't)

B15		fx									
	A	B	C	D	E	F	G	H	I	J	K
1											
2			F	P1	P2	P3	P4	P5			
3		Solution	0	0	0	0	0	0			
4		9-10a	1	1					= \$C\$3*C4+\$D\$3*D4+\$E\$3*E4+\$F\$3*F4+\$G\$3*G4+\$H\$3*H4	10	>=
5		10-11a	1	1	1				= \$C\$3*C5+\$D\$3*D5+\$E\$3*E5+\$F\$3*F5+\$G\$3*G5+\$H\$3*H5	12	>=
6		11a - Noon	0.5	1	1	1			= \$C\$3*C6+\$D\$3*D6+\$E\$3*E6+\$F\$3*F6+\$G\$3*G6+\$H\$3*H6	14	>=
7		Noon - 1p	0.5	1	1	1	1		= \$C\$3*C7+\$D\$3*D7+\$E\$3*E7+\$F\$3*F7+\$G\$3*G7+\$H\$3*H7	16	>=
8		1 - 2p	1		1	1	1	1	= \$C\$3*C8+\$D\$3*D8+\$E\$3*E8+\$F\$3*F8+\$G\$3*G8+\$H\$3*H8	18	>=
9		2-3p	1			1	1	1	= \$C\$3*C9+\$D\$3*D9+\$E\$3*E9+\$F\$3*F9+\$G\$3*G9+\$H\$3*H9	17	>=
10		3-4p	1				1	1	= \$C\$3*C10+\$D\$3*D10+\$E\$3*E10+\$F\$3*F10+\$G\$3*G10+\$H\$3*H10	15	>=
11		4-5p	1					1	= \$C\$3*C11+\$D\$3*D11+\$E\$3*E11+\$F\$3*F11+\$G\$3*G11+\$H\$3*H11	10	>=
12		Full Time	1						= \$C\$3*C12+\$D\$3*D12+\$E\$3*E12+\$F\$3*F12+\$G\$3*G12+\$H\$3*H12	12	<=
13		Part Time		4	4	4	4	4	= \$C\$3*C13+\$D\$3*D13+\$E\$3*E13+\$F\$3*F13+\$G\$3*G13+\$H\$3*H13	56	<=
14		Objective	100	32	32	32	32	32	= \$C\$3*C14+\$D\$3*D14+\$E\$3*E14+\$F\$3*F14+\$G\$3*G14+\$H\$3*H14		Min
15											
16											

Labor Planning (con't)

Excel spreadsheet showing labor planning data and the Solver Parameters dialog box.

Excel Spreadsheet Data:

	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O
1															
2			F	P1	P2	P3	P4	P5							
3		Solution	0	0	0	0	0	0							
4		9-10a	1	1						0	10	>=			
5		10-11a	1	1	1					0	12	>=			
6		11a - Noon	0.5	1	1	1				0	14	>=			
7		Noon - 1p	0.5	1	1	1	1			0	16	>=			
8		1 - 2p	1		1	1	1	1		0	18	>=			
9		2-3p	1			1	1	1		0	17	>=			
10		3-4p	1				1	1		0	15	>=			
11		4-5p	1					1		0	10	>=			
12		Full Time	1							0	12	<=			
13		Part Time		4	4	4	4	4		0	56	<=			
14		Objective	100	32	32	32	32	32	0		Min				

Solver Parameters Dialog Box:

- Set Target Cell:
- Equal To: ☐ Max ☒ Min ☐ Value of:
- By Changing Cells:
- Subject to the Constraints:
 -
 -
 -
 -
 -
 -

Buttons: Solve, Close, Options, Add, Change, Delete, Reset All, Help.

Labor Planning (con't)

	B15																
	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q
1																	
2			F	P1	P2	P3	P4	P5									
3		Solution	10	3	3	3	2	3									
4		9-10a	1	1					13	10	>=						
5		10-11a	1	1	1				16	12	>=						
6		11a - Noon	0.5	1	1	1			14	14	>=						
7		Noon - 1p	0.5	1	1	1	1		16	16	>=						
8		1 - 2p	1		1	1	1	1	21	18	>=						
9		2-3p	1			1	1	1	18	17	>=						
10		3-4p	1				1	1	15	15	>=						
11		4-5p	1					1	13	10	>=						
12		Full Time	1						10	12	<=						
13		Part Time		4	4	4	4	4	56	56	<=						
14		Objective	100	32	32	32	32	32	1448		Min						
15																	
16																	
17																	
18																	
19																	
20																	
21																	
22																	
23																	
24																	
25																	

Solver Results

Solver found a solution. All constraints and optimality conditions are satisfied.

☒ Keep Solver Solution
 ☐ Restore Original Values

Reports

Answer
 Sensitivity
 Limits

OK

Cancel

Save Scenario...

Help

Labor Planning (con't)

Another solution at \$1448 – there are multiple optimal solutions

	A	B	C	D	E	F	G	H	I	J	K	
1												
2			F	P1	P2	P3	P4	P5				
3		Solution	10	0	2	7	5	0				
4		9-10a	1	1					10	10	>=	
5		10-11a	1	1	1				12	12	>=	
6		11a - Noon	0.5	1	1	1			14	14	>=	
7		Noon - 1p	0.5	1	1	1	1		19	16	>=	
8		1 - 2p	1		1	1	1	1	24	18	>=	
9		2-3p	1			1	1	1	22	17	>=	
10		3-4p	1				1	1	15	15	>=	
11		4-5p	1					1	10	10	>=	
12		Full Time	1						10	12	<=	
13		Part Time		4	4	4	4	4	56	56	<=	
14		Objective	100	32	32	32	32	32	1448		Min	
15												
16												

Labor Planning (con't)

Another Optimal solution – which one depends upon starting values for the solution variables, and the order of the constraints !!!

N7 fx												
	A	B	C	D	E	F	G	H	I	J	K	L
1												
2			F	P1	P2	P3	P4	P5				
3		Solution	10	6	1	2	5	0				
4		9-10a	1	1					16	10	>=	
5		10-11a	1	1	1				17	12	>=	
6		11a - Noon	0.5	1	1	1			14	14	>=	
7		Noon - 1p	0.5	1	1	1	1		19	16	>=	
8		1 - 2p	1		1	1	1	1	18	18	>=	
9		2-3p	1			1	1	1	17	17	>=	
10		3-4p	1				1	1	15	15	>=	
11		4-5p	1					1	10	10	>=	
12		Full Time	1						10	12	<=	
13		Part Time		4	4	4	4	4	56	56	<=	
14		Objective	100	32	32	32	32	32	1448		Min	
15												

Mixed-Integer Programming Problem Example

- There are many situations in which **some of the variables** are restricted to be integers and some are not
- As an example consider a chemical company that produces two industrial chemicals
- Xylene must be produced in 50-pound bags
- Hexall is sold by the pound and can be produced in any quantity
- Both xylene and hexall are composed of three ingredients – A , B , and C
- Xylene sells for \$85 a bag and hexall for \$1.50 per pound

Mixed-Integer Programming Problem Example (con't)

AMOUNT PER 50-POUND BAG OF XYLINE (LB)	AMOUNT PER POUND OF HEXALL (LB)	AMOUNT OF INGREDIENTS AVAILABLE
30	0.5	2,000 lb–ingredient A
18	0.4	800 lb–ingredient B
2	0.1	200 lb–ingredient C

- We want to maximize profit
- We let X = number of 50-pound bags of xyline
- We let Y = number of pounds of hexall
- This is a mixed-integer programming problem as Y is not required to be an integer

Mixed-Integer Programming Problem Example (con't)

■ The model is

Maximize profit = $\$85X + \$1.50Y$

subject to $30X + 0.5Y \leq 2,000$

$18X + 0.4Y \leq 800$

$2X + 0.1Y \leq 200$

$X, Y \geq 0$ and X integer

■ Using Excel QM

[illegible]

Mixed-Integer Programming Problem Example (con't)

G10 Σ =SUMPRODUCT(B10:C10,\$B\$16:\$C\$16)

Linear Programming

Enter the values in the shaded area. Then go to the DATA Tab on the ribbon, click on Solver in the Data Analysis Group and then click SOLVE.
If SOLVER is not on the Data Tab then please see the Help file (Solver) for instructions.

Signs

<	less than or equal to
=	equals (You need to enter an apostrophe first.)
>	greater than or equal to

Data

	x 1	x 2	sign	RHS
Objective	85	1.5		
Constraint 1	30	0.5	<	2000
Constraint 2	18	0.4	<	800
Constraint 3	2	0.1	<	200

Results

LHS	Slack/Surplus
3770	
1330	670
800	0
90	110

Problem setup area

< constraints	> constraints		
1330	2000	0	0
800	800	0	0
90	200	0	0

Results

Variables			
44	20		
Objective			3770

Solver Parameters

Set Target Cell: G10

Equal To: ☒ Max ☐ Min ☐ Value of: 0

By Changing Cells: B16:C16

Subject to the Constraints:

- B16:C16 = integer
- J11:J13 <= K11:K13
- L11:L13 >= M11:M13

Need to manually Add integer constraint

Real World Mixed Integer Problem

[Network Mode Optimization for the DHL Supply Chain, *Informatics J. on Applied Analytics*, Vol 51, No 3 2021]

Objective function

$$\text{Min} \sum_{k \in V} \left\{ f \cdot \theta_k + \sum_{i \in I_0} \sum_{j \in I_0} x_{ijk} \cdot (c_{ij} + p) - \sum_{i \in I_0} x_{i0k} \cdot p \right\} + \sum_{i \in I} \lambda_i \cdot y_i$$

Subject to:

Degree constraints

$$\sum_{i \in I_0} x_{ijk} = \sum_{h \in I_0} x_{jhk}, \forall k \in V, \forall j \in I_0$$

$$x_{iik} = 0, \forall i \in I_0, \forall k \in V$$

$$\sum_{k \in V} \sum_{j \in I_0} x_{ijk} + y_i = 1, \forall i \in I$$

$$\sum_{j \in I_0} x_{0jk} \leq 1, \forall k \in V$$

Time-window constraints

$$a_j \cdot x_{ijk} \geq s_{ijk} \geq e_j \cdot x_{ijk}, \forall i \in I_0, \forall j \in I, \forall k \in V$$

$$s_{ijk} \leq \sum_{h \in I_0} s_{hik} + x_{ijk} \cdot \left(\frac{d_{ij}}{g} + \mu_i \right) + b \cdot l_{jk}, \forall i \in I, \forall j \in I_0, \forall k \in V$$

$$s_{ijk} \geq \sum_{h \in I_0} s_{hik} + x_{ijk} \cdot \left(\frac{d_{ij}}{g} + \mu_i \right) + b \cdot l_{jk} - (1 - x_{ijk}) \cdot M, \forall i \in I, \forall j \in I_0, \forall k \in V$$

Layover constraints

$$\sum_{i \in I_0} \sum_{j \in I_0} x_{ijk} \cdot \left(\frac{d_{ij}}{g} + \mu_i \right) \leq \delta + 2 \cdot \delta \cdot z_k, \forall k \in V$$

$$\sum_{i \in I_0} \sum_{j \in I_0} x_{ijk} \cdot \left(\frac{d_{ij}}{g} + \mu_i \right) \leq 2 \cdot \delta + \delta \cdot r_k, \forall k \in V$$

$$r_k + z_k = \sum_{j \in I_0} l_{jk}, \forall k \in V$$

$$l_{jk} \leq \sum_{i \in I_0} x_{ijk}, \forall j \in I_0, \forall k \in V$$

Truck-capacity constraints

$$q + (w_i - q) \cdot \sum_{k \in V} x_{0ik} \geq u_i \geq w_i, \forall i \in I$$

$$u_i - u_j + q \cdot \left(\sum_{k \in V} x_{ijk} \right) \leq q - w_j, \forall i, j \in I$$

Maximum-travel-time constraints

$$t_k^{\max} \geq t_k \geq \sum_{i \in I_0} s_{i0k} - \sum_{j \in I_0} s_{0jk} + \sum_{j \in I_0} x_{0jk} \cdot \left(\frac{d_{0j}}{g} \right), \forall k \in V$$

$$\theta_k \geq \frac{t_k}{2\Delta}, \forall k \in V$$

Intranode-distance constraints

$$d_{ij} \cdot x_{ijk} \leq d^{\max}, \forall i, j \in I_0, \forall k \in V$$

Symmetry breaking inequalities

$$t_k \geq t_{k+1}, \forall k \in V$$

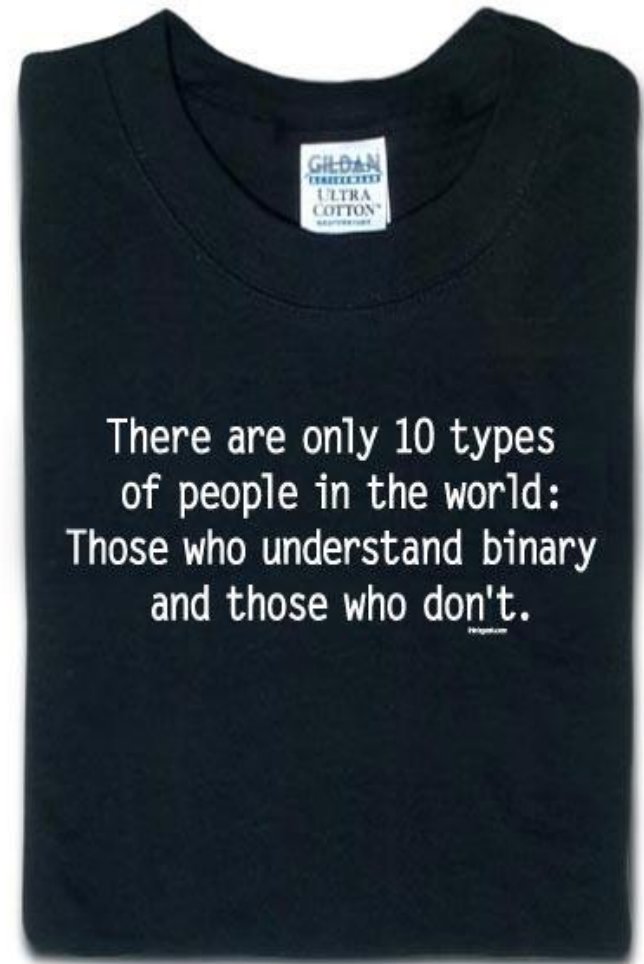
Integrality and nonnegativity constraints

$$x_{ijk}, y_i, l_{ik}, r_k, z_k \in \{0, 1\}, y_0 = 0, \forall i \in I_0, \forall j \in I_0, \forall k \in V$$

$$t_k \geq 0, \theta_k \in \mathbb{Z}_+, s_{ijk} \geq 0, u_i \geq 0, \forall i \in I, \forall j \in I_0, \forall k \in V$$

Modeling With 0-1 (Binary) Variables

- Let's demonstrate how 0-1 variables can be used to model several diverse situations
- Typically a 0-1 variable is assigned a value of 0 if a certain condition is not met and a 1 if the condition is met
- This is also called a *binary variable*

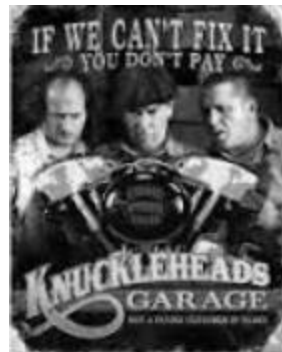


The Assignment Problem

- Another common LP algorithm is the assignment method
- Each assignment problem has associated with it a table, or matrix
- Generally, the rows contain the objects or people we wish to assign, and the columns comprise the tasks or things to which we want them assigned
- The numbers in the table are the costs associated with each particular assignment
- An assignment problem can be viewed as a transportation problem in which the capacity from each source is 1 and the demand at each destination is 1

Assignment Model Approach

- The Knucklehead's Fix-It Shop has three rush projects to repair
- The shop has three repair persons with different talents and abilities (Larry, Curley, and Moe)
- The owner has estimates of wage costs for each worker for each project
- The owner's objective is to assign the three projects to the workers in a way that will result in the lowest cost to the shop
- Each project will be assigned exclusively to one worker



Assignment Problem Approach (con't)

Estimated Project Repair Costs for the Fix-It Shop
Assignment Problem

PERSON	PROJECT		
	1	2	3
Larry	\$11	\$14	\$6
Curley	8	10	11
Moe	9	12	7





PERSON	PROJECT		
	1	2	3
Larry	\$11	\$14	\$6
Curley	8	10	11
Moe	9	12	7

- $X_{ij} = 1$ if person i is assigned to job j , else 0
- Minimize $Z = 11X_{11} + 14X_{12} + 6X_{13} +$
- $8X_{21} + 10X_{22} + 11X_{23} +$
- $9X_{31} + 12X_{32} + 7X_{33}$
- Subject to:
 - $X_{i1} + X_{i2} + X_{i3} \leq 1$ for $i = 1$ to 3
 - $X_{ij} \geq 0$ and $X_{ij} \leq 1$ and X_{ij} is an integer (X_{ij} are binary)

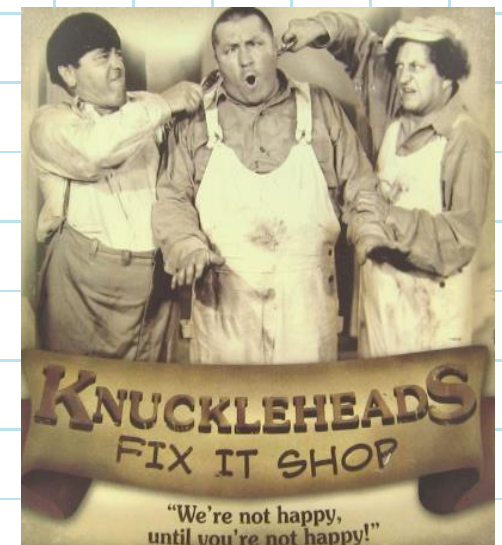
Assignment Model Approach (con't)

B	C	D	E	F
Cost for Assignments				
Project 1	Project 2	Project 3		
11	14	6		
8	10	11		
9	12	7		
Made				
Project 1	Project 2	Project 3	Total projects	Supply
0	0	1	1	1
0	1	0	1	1
1	0	0	1	1
1	1	1		
1	1	1		
25				

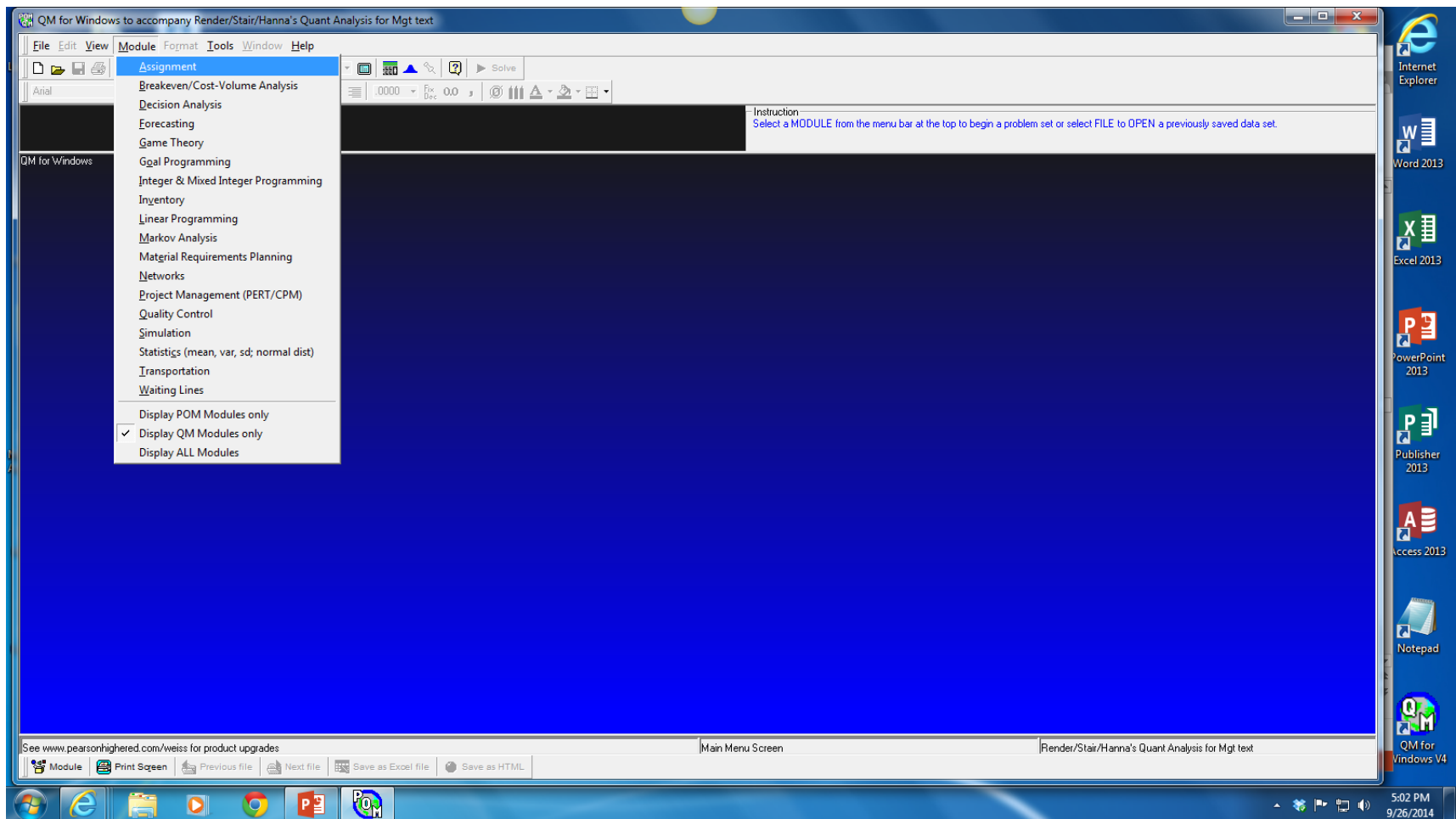
E
10 =SUM(B10:D10)

B
13 =SUM(B10:B12)

B
16 =SUMPRODUCT(B3:D5,B10:D12)



Using QM...



Using QM...

Create data set for Assignment

Title: KnuckleHead's Shop Modify default title

Number of Jobs
3

Number of Machines
3

Objective
☐ Maximize
☒ Minimize

Row names Column names Overview

☒ Job 1, Job 2, Job 3, ...
☐ a, b, c, d, e, ...
☐ A, B, C, D, E, ...
☐ 1, 2, 3, 4, 5, ...
☐ January, February, March, April, ...
Click here to set start month
☐ Other

Cancel Help OK

Using QM...

QM for Windows - [Data Table]

File Edit View Module Format Tools Window Help

100% Solve

Arial 8.25 B I U .00 Fix Dec 0.0

Objective
☐ Maximize
☒ Minimize

Instruction
Enter the cost of assigning moe to project

	Project 1	Project 2	Project 3
Larry	11	14	6
Curley	8	10	11
Moe	9	12	7

Using QM...

QM for Windows

File Edit View Module Format Tools Window Help

100%

8.25

Objective

☐ Maximize

☒ Minimize

Instruction

There are more results available in additional windows. These may be opened by using the WINDOW of

Assignments

KnuckleHead's Shop Solution

	Project 1	Project 2	Project 3
Optimal cost = \$25			
Larry	11	14	Assign 6
Curley	8	Assign 10	11
Moe	Assign 9	12	7

Assignment List

KnuckleHead's Shop Solution

JOB	Assigned to	Cost
Larry	Project 3	6
Curley	Project 2	10
Moe	Project 1	9
Total		25

Fixed-Charge Problem Example

- Often businesses are faced with decisions involving a fixed charge that will affect the cost of future operations
- A manufacturing company is planning to build **at least one** new plant and three cities are being considered:
 - Baytown, Texas
 - Lake Charles, Louisiana
 - Mobile, Alabama
- Once the plant or plants are built, the company wants to have capacity to produce at least 38,000 units each year

Fixed-Charge Problem Example (con't)

- Fixed and variable costs for Manufacturing Co.

SITE	ANNUAL FIXED COST	VARIABLE COST PER UNIT	ANNUAL CAPACITY
Baytown, TX	\$340,000	\$32	21,000
Lake Charles, LA	\$270,000	\$33	20,000
Mobile, AL	\$290,000	\$30	19,000

SITE	ANNUAL FIXED COST	VARIABLE COST PER UNIT	ANNUAL CAPACITY
Baytown, TX	\$340,000	\$32	21,000
Lake Charles, LA	\$270,000	\$33	20,000
Mobile, AL	\$290,000	\$30	19,000

- Management decisions ? [need 38,000/yr]
- What are the variables ?
- Use of binary variables?



■ Do not look ahead !



Fixed-Charge Problem Example (con't)

- We can define the decision variables as

$$X_1 = \begin{cases} 1 & \text{if factory is built in Baytown} \\ 0 & \text{otherwise} \end{cases}$$

$$X_2 = \begin{cases} 1 & \text{factory is built in Lake Charles} \\ 0 & \text{otherwise} \end{cases}$$

$$X_3 = \begin{cases} 1 & \text{if factory is built in Mobile} \\ 0 & \text{otherwise} \end{cases}$$

X_4 = number of units produced at Baytown plant

X_5 = number of units produced at Lake Charles plant

X_6 = number of units produced at Mobile plant

Fixed-Charge Problem Example (con't)

- The integer programming formulation becomes

$$\begin{aligned}\text{Minimize cost} = & 340,000X_1 + 270,000X_2 + 290,000X_3 \\ & + 32X_4 + 33X_5 + 30X_6\end{aligned}$$

$$\begin{aligned}\text{subject to} \quad & X_4 + X_5 + X_6 \geq 38,000 \\ & X_4 \leq 21,000X_1 \\ & X_5 \leq 20,000X_2 \\ & X_6 \leq 19,000X_3 \\ & X_1, X_2, X_3 = 0 \text{ or } 1; \\ & X_4, X_5, X_6 \geq 0 \text{ and integer}\end{aligned}$$

Note capacity constraints
expressed in terms of
binary variables

- The optimal solution is

$$X_1 = 0, X_2 = 1, X_3 = 1 \text{ [build Lake Charles \& Mobile Plants]}$$

$$X_4 = 0, X_5 = 19,000, X_6 = 19,000$$

$$\text{Objective function value} = \$1,757,000$$

Fixed-Charge Problem Example (con't)

FIXED-CHARGE.xls - Microsoft Excel

Home Insert Page Layout Formulas Data Review View Developer

Get External Data Refresh All Edit Links Connections Sort & Filter Filter Clear Reapply Advanced Data Validation Consolidate What-If Analysis Text to Columns Remove Duplicates Outline Analysis

C18

	A	B	C	D	E	F	G	H	I	
1										
2		X1	X2	X3	X4	X5	X6		RHS	
3	Solution	0	0	0	0	0	0			
4	Con 1				1	1	1	=B\$3*B4+C\$3*C4+D\$3*D4+E\$3*E4+F\$3*F4+G\$3*G4	38000	>=
5	Con 2	-28000			1			=B\$3*B5+C\$3*C5+D\$3*D5+E\$3*E5+F\$3*F5+G\$3*G5	0	<=
6	Con 3		-20000			1		=B\$3*B6+C\$3*C6+D\$3*D6+E\$3*E6+F\$3*F6+G\$3*G6	0	<=
7	Con 4			-19000			1	=B\$3*B7+C\$3*C7+D\$3*D7+E\$3*E7+F\$3*F7+G\$3*G7	0	<=
8	Objective	340000	270000	290000	32	33	30	=B\$3*B8+C\$3*C8+D\$3*D8+E\$3*E8+F\$3*F8+G\$3*G8		Min
9										
10										
11										
12										
13										

Fixed-Charge Problem Example (con't)

H8 f_x $=\$B\$3*B8+\$C\$3*C8+\$D\$3*D8+\$E\$3*E8+\$F\$3*F8+\$G\$3*G8$									
	A	B	C	D	E	F	G	H	I
1									
2		X1	X2	X3	X4	X5	X6		RHS
3	Solution	0	0	0	0	0	0		
4	Con 1				1	1	1	$=\$B\$3*B4+\$C\$3*C4+\$D\$3*D4+\$E\$3*E4+\$F\$3*F4+\$G\$3*G4$	38000 \geq
5	Con 2	-28000			1			$=\$B\$3*B5+\$C\$3*C5+\$D\$3*D5+\$E\$3*E5+\$F\$3*F5+\$G\$3*G5$	0 \leq
6	Con 3		-20000			1		$=\$B\$3*B6+\$C\$3*C6+\$D\$3*D6+\$E\$3*E6+\$F\$3*F6+\$G\$3*G6$	0 \leq
7	Con 4			-19000			1	$=\$B\$3*B7+\$C\$3*C7+\$D\$3*D7+\$E\$3*E7+\$F\$3*F7+\$G\$3*G7$	0 \leq
8	Objective	340000	270000	290000	32	33	30	$=\$B\$3*B8+\$C\$3*C8+\$D\$3*D8+\$E\$3*E8+\$F\$3*F8+\$G\$3*G8$	Min

Solver Parameters

Set Target Cell:

Equal To: ☐ Max ☒ Min ☐ Value of:

By Changing Cells:

Subject to the Constraints:

- $\$B\$3:\$D\$3 = \text{binary}$
- $\$E\$3:\$G\$3 = \text{integer}$
- $\$E\$3:\$G\$3 \geq 0$
- $\$H\$4 \geq \$I\4
- $\$H\$5 \leq \$I\5
- $\$H\$6 \leq \$I\6

Fixed-Charge Problem Example (con't)

FIXED-CHARGE.xls - Microsoft Excel

Home Insert Page Layout Formulas Data Review View Developer

Get External Data Refresh All Edit Links Connections Sort & Filter Filter Sort Advanced Clear Reapply Data Tools Text to Columns Remove Duplicates Data Validation Consolidate What-If Analysis Group Ungroup Subtotal Data Analysis Solver

C18

	A	B	C	D	E	F	G	H	
1									
2		X1	X2	X3	X4	X5	X6		
3	Solution	0	1	1	4.47286652160983E-12	18999.9999999977	19000		
4	Con 1				1	1		=B\$3*B4+C\$3*C4+D\$3*D4+E\$3*E4+F\$3*F4+G\$3*G4	38000
5	Con 2	-28000			1			=B\$3*B5+C\$3*C5+D\$3*D5+E\$3*E5+F\$3*F5+G\$3*G5	0
6	Con 3		-20000			1		=B\$3*B6+C\$3*C6+D\$3*D6+E\$3*E6+F\$3*F6+G\$3*G6	0
7	Con 4			-19000			1	=B\$3*B7+C\$3*C7+D\$3*D7+E\$3*E7+F\$3*F7+G\$3*G7	0
8	Objective	340000	270000	290000	32	33	30	=B\$3*B8+C\$3*C8+D\$3*D8+E\$3*E8+F\$3*F8+G\$3*G8	
9									
10									
11									
12									
13									
14									
15									

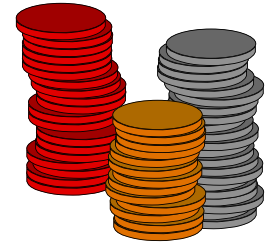
Portfolio Optimization

- Portfolio problems can be regular LP problems or a type of integer problems
 - Maximize the return from a number of financial investments subject to risk and diversity constraints – regular LP
 - Maximize the return on projects by selecting which projects to work on within risk and budget constraints – a zero/one type of integer programming problem (binary programming)

Net Present Value

- $NPV = \sum (B - C)_t / (1+i)^t$
- Where $(B-C)_t$ is the benefit minus the cost for period t
- i is the interest rate (cost of borrowing money or opportunity cost for other uses of cash)
- For NPV, benefit minus cost is more formally revenue (cash in) minus expenditures (cash out)

NPV Example

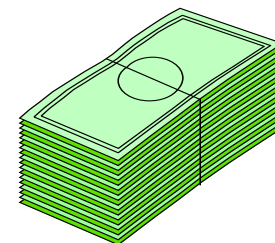


Year	Benefit	Cost	B-C	Discounted B-C
1	\$0.00	\$175,000.00	-\$175,000.00	-\$159,090.91
2	\$0.00	\$175,000.00	-\$175,000.00	-\$144,628.10
3	\$50,000.00	\$25,000.00	\$25,000.00	\$18,782.87
4	\$100,000.00	\$10,000.00	\$90,000.00	\$61,471.21
5	\$100,000.00	\$10,000.00	\$90,000.00	\$55,882.92
6	\$100,000.00	\$10,000.00	\$90,000.00	\$50,802.65
7	\$100,000.00	\$10,000.00	\$90,000.00	\$46,184.23
8	\$100,000.00	\$10,000.00	\$90,000.00	\$41,985.66
9	\$100,000.00	\$10,000.00	\$90,000.00	\$38,168.79
10	\$100,000.00	\$10,000.00	\$90,000.00	\$34,698.90
Total	\$750,000.00	\$445,000.00	\$305,000.00	\$44,258.22
		Interest = 0.1		

Internal Rate of Return (IRR)

- Another similar project financial evaluation technique is called the ***internal rate of return*** (IRR)
- This metric is better than NPV since it is not as sensitive to the uncertainties of future benefits and costs and to the future interest rates
- The internal rate of return is the value of the interest rate that yields a zero value for NPV
 - This can be calculated in spreadsheet programs by using built-in “solver” tools. Since in reality a quadratic equation is being solved, multiple IRR values could be found. Thus one must impose additional constraints on the solution (such as IRR is positive, or in a given range).

IRR Example



Year	Benefit	Cost	B-C	Discounted B-C
1	\$0.00	\$175,000.00	-\$175,000.00	-\$154,728.74
2	\$0.00	\$175,000.00	-\$175,000.00	-\$136,805.61
3	\$50,000.00	\$25,000.00	\$25,000.00	\$17,279.80
4	\$100,000.00	\$10,000.00	\$90,000.00	\$55,001.46
5	\$100,000.00	\$10,000.00	\$90,000.00	\$48,630.33
6	\$100,000.00	\$10,000.00	\$90,000.00	\$42,997.19
7	\$100,000.00	\$10,000.00	\$90,000.00	\$38,016.58
8	\$100,000.00	\$10,000.00	\$90,000.00	\$33,612.90
9	\$100,000.00	\$10,000.00	\$90,000.00	\$29,719.32
10	\$100,000.00	\$10,000.00	\$90,000.00	\$26,276.76
Total	\$750,000.00	\$445,000.00	\$305,000.00	\$0.00
		Interest = 0.131011619		

Projects with the same net present value may have different internal rates of return

Case 1

Period	Benefit	Cost	B-C	Discounted B-C
1	0	70	-70	-\$60.87
2	0	50	-50	-\$37.81
3	20	30	-10	-\$6.58
4	90	0	90	\$51.46
5	120	0	120	\$59.66
			NPV:	\$5.87
		Interest:	0.15	
		IRR:	0.17	

Case 2

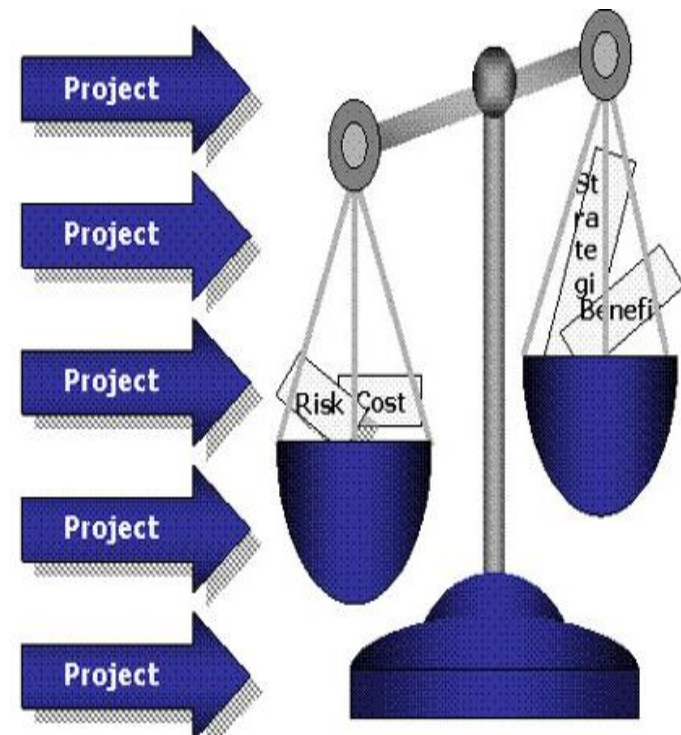
Period	Benefit	Cost	B-C	Discounted B-C
1	0	20	-20	-\$17.39
2	0	40	-40	-\$30.25
3	20	50	-30	-\$19.73
4	90	55	35	\$20.01
5	120	12.95	107.05	\$53.22
			NPV:	\$5.87
		Interest:	0.15	
		IRR:	0.19	

Example Project Portfolio Data

	B	C	D	E
	Project	IRR	Risk Factor	Cost (\$)
	1	0.35	0.15	200
	2	0.25	0.05	500
	3	0.3	0.5	700
	4	0.15	0.3	300
	5	0.28	0.25	400
	6	0.25	0.3	900
	7	0.2	0.2	600
	8	0.3	0.1	800

Constraints

- Must balance risk and reward
- Cannot exceed budget of \$3000
- Risk must be within limit of “management reserve” (percentage of overall budget, 20% in this example - \$600)



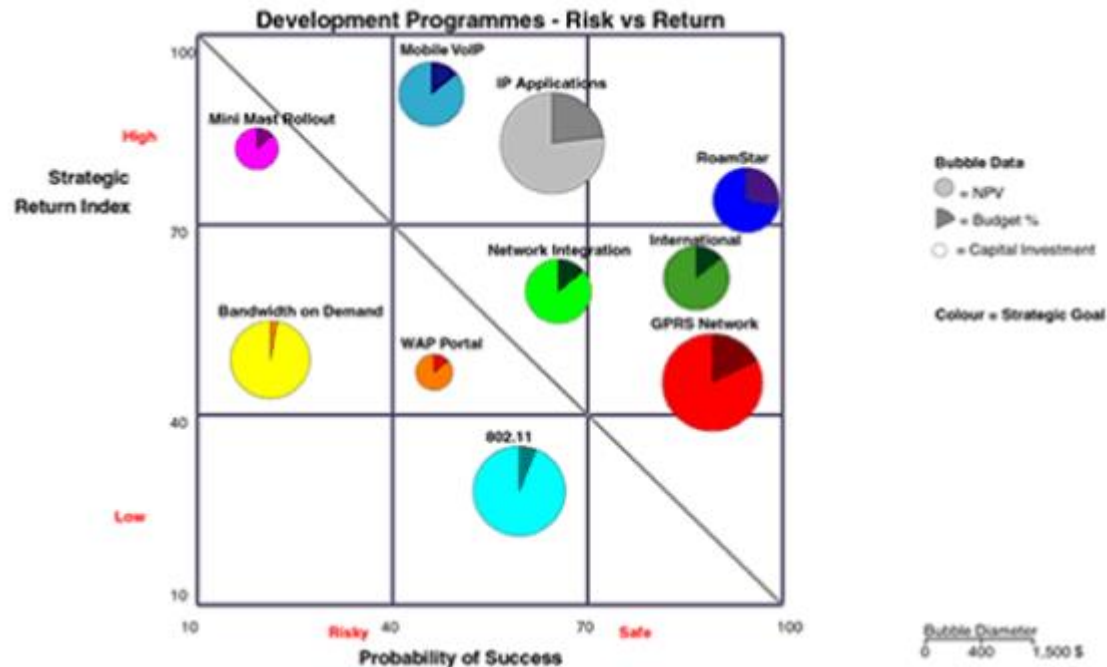
Which projects would you do ?

Budget less than \$3000

Risk Contingency \leq \$600

	B	C	D	E
	Project	IRR	Risk Factor	Cost (\$)
	1	0.35	0.15	200
	2	0.25	0.05	500
	3	0.3	0.5	700
	4	0.15	0.3	300
	5	0.28	0.25	400
	6	0.25	0.3	900
	7	0.2	0.2	600
	8	0.3	0.1	800

- What are the variables ?
- Use of binary variables?



■ Do not look ahead !



[column E is a binary value]

[illegible]

Using Excel Solver to Set Up Constraints and Get Solution

[newer versions of Excel have “bin” constraint]

	B	C	D	E	F	G	H	I
	Project	IRR	Risk Factor	Cost (\$)	Invest(%)	Invest (\$)	Return (\$)	Risk
	1	0.35	0.15	200	0	0	0	0
	2	0.25	0.05	500	0	0	0	0
	3	0.3	0.5	700	0	0	0	0
	4	0.15	0.3	300	0	0	0	0
	5	0.28	0.25	400	0	0	0	0
	6	0.25	0.3	900	0	0	0	0
	7	0.2	0.2	600	0	0	0	0
	8	0.3	0.1	800	0	0	0	0
						0	0	0
						Tot Invest	Tot Return	Tot Risk

Solver Parameters

Set Target Cell:

Equal To: ☒ Max ☐ Min ☐ Value of:

By Changing Cells:

Subject to the Constraints:

Answers

[not doing projects 3,4,5]

	B	C	D	E	F	G	H	I	
	Project	IRR	Risk Factor	Cost (\$)	Invest(%)	Invest (\$)	Return (\$)	Risk	
	1	0.35	0.15	200	1	200	70	30	
	2	0.25	0.05	500	1	500	125	25	
	3	0.3	0.5	700	0	0	0	0	
	4	0.15	0.3	300	0	0	0	0	
	5	0.28	0.25	400	0	0	0	0	
	6	0.25	0.3	900	1	900	225	270	
	7	0.2	0.2	600	1	600	120	120	
	8	0.3	0.1	800	1	800	240	80	
						3000	780	525	
						Tot Invest	Tot Return	Tot Risk	

Limits on Alternatives Selected

- Suppose it is required to select no more than Z of the projects *regardless* of the funds available

- This would require adding a constraint

$$X_1 + X_2 + X_3 + \dots \leq Z$$

- If they had to fund *exactly* Z of the projects the constraint would be

$$X_1 + X_2 + X_3 + \dots = Z$$

Dependent Selections

- At times the selection of one project depends on the selection of another project
- Suppose project 1 could only be done if the project 2 was also done
- The following constrain would force this to occur

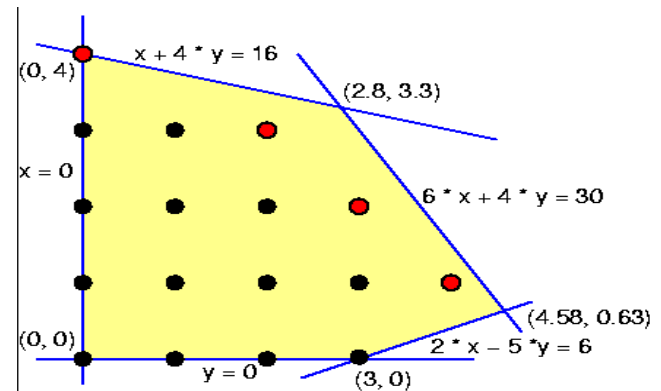
$$X_1 \leq X_2 \quad \text{or} \quad X_1 - X_2 \leq 0$$

- If we wished for project 1 and project 2 to either both be selected or both not be selected, the constraint would be

$$X_1 = X_2 \quad \text{or} \quad X_1 - X_2 = 0$$

Homework

- Textbook Chapter 9 thru section 9.4
- Discussion questions from chapter 9: 2, 3, 4
- Project Seven →



Project 7



- A aero company uses three plants to manufacture amphibious airplane floats
- Let x_1 be the batches of big float batches per week and let x_2 be the number of small float batches per week
- The big floats contribute 3 units to profit per batch and the small floats contribute 2 units to profit per batch



Project 7 (con't)

- Plant 1 can produce up to 4 batches of x_1 per week; plant 2 can produce up to 6 units of x_2 per week; plant 3 can produce up to 18 batches per week in the ratio of 3 of x_1 and 2 of x_2
- What is the optimal mix of big float and small float batches to make each week?
- The number of float batches must be integers
- Show your setup and solution in Excel or QM

